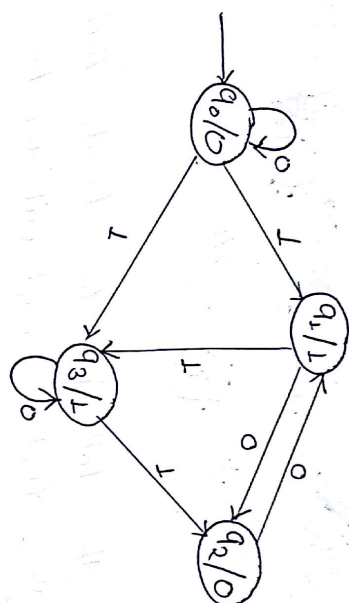


Theory of Computation Branch :- Computer Science (M sem)
 Q.1. Construct a mealy machine equivalent to the given Moore machine.



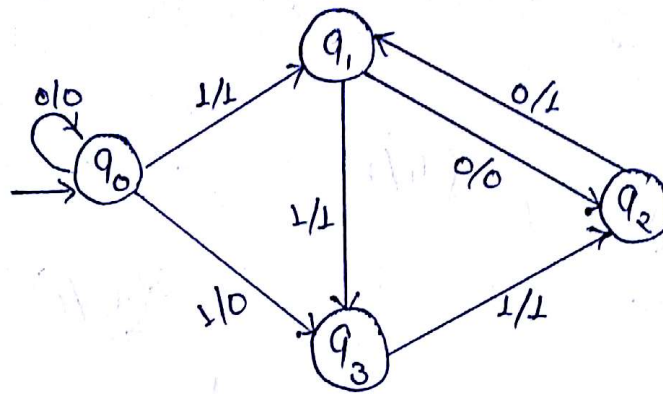
Sol. Step 1:

Convert a moore diagram into a mealy transition table.

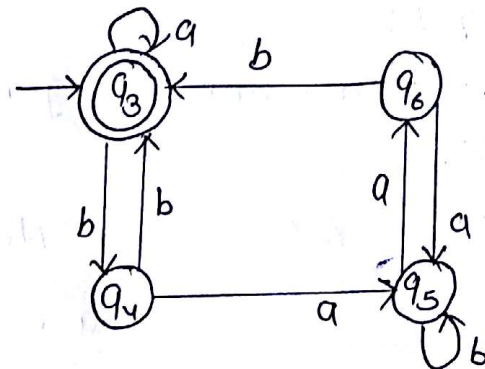
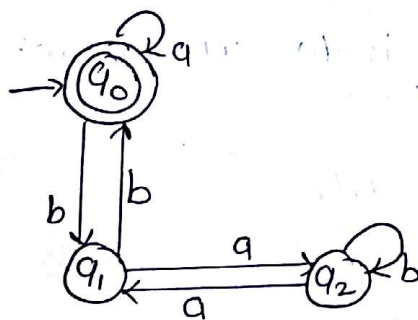
Present state	Next state		output
	a=0	a=1	
→ q ₀	q ₀	q ₁	0/P
q ₁	q ₂	q ₃	1
q ₂	q ₁	q ₃	0
q ₃	q ₃	q ₀	1

Step 2: state transition table for mealy machine.

Present state	Next state			
	a=0		a=1	
	State	output	State	output
→ q ₀	q ₀	0	q ₁	1
q ₁	q ₂	0	q ₃	1
q ₂	q ₁	1	q ₃	1
q ₃	q ₃	1	q ₀	0



Q2). Consider the following two finite Automata FA1 & FA2 over $\{0,1\}$. Find out whether they are equivalent or not.



Sol. Comparison table :-

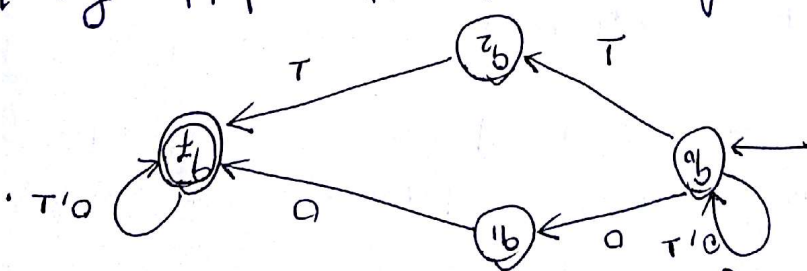
State	For Input a	For Input b
(q, q')	(q_a, q'_a)	(q_b, q'_b)
(q_0, q_3)	(q_0, q_3)	(q_1, q_4)
(q_1, q_4)	(q_2, q_5)	(q_0, q_3)
(q_2, q_5)	(q_1, q_6)	(q_2, q_5)
(q_1, q_6)	(q_2, q_5)	(q_0, q_3)

In each entry in table (q, q') is such q is final state at FA1 then q' is also final state else both are non-final.

This FA1 & FA2 are equivalent.

Theory of Computation

Convert the following NFA system to DFA.



Soln: - Construction of state transition table for the given NFA.

States	0	1
q0	q1	q2
q1	qf	qf
q2	qf	qf
qf	qf	qf

Step 1: - Input symbol for the resultant DFA = Input symbols for the given NFA

Step 2: - Start state for the resultant DFA = 0 & 1.

Step 3: - Since there are 4 states in the given NFA, so the no. of states in the resultant DFA = 2⁴.

All the states of resultant DFA = set of subset of the set of states of the given NFA

Step 4: - we will find out the accessible states from the step 3.

$$\delta(q_0, 0) = [q_0, q_1] \quad \delta(q_0, 1) = [q_0, q_2]$$

$$\delta(q_1, 0) = [q_1] \quad \delta(q_1, 1) = [q_1, q_f]$$

$$\delta(q_2, 0) = [q_2] \quad \delta(q_2, 1) = [q_2, q_f]$$

$$\delta(q_f, 0) = [q_f] \quad \delta(q_f, 1) = [q_f]$$

$$= [q_0, q_2] \cup \phi$$

$$\delta(q_0, q_1) = \delta(q_0, 1) \cup \delta(q_1, 1)$$

$$= [q_0, q_1, q_2]$$

$$= [q_0, q_1] \cup [q_2]$$

$$\delta(q_0, q_1, q_2) = \delta(q_0, 0) \cup \delta(q_1, 0) \cup \delta(q_2, 0)$$

The accessible states are

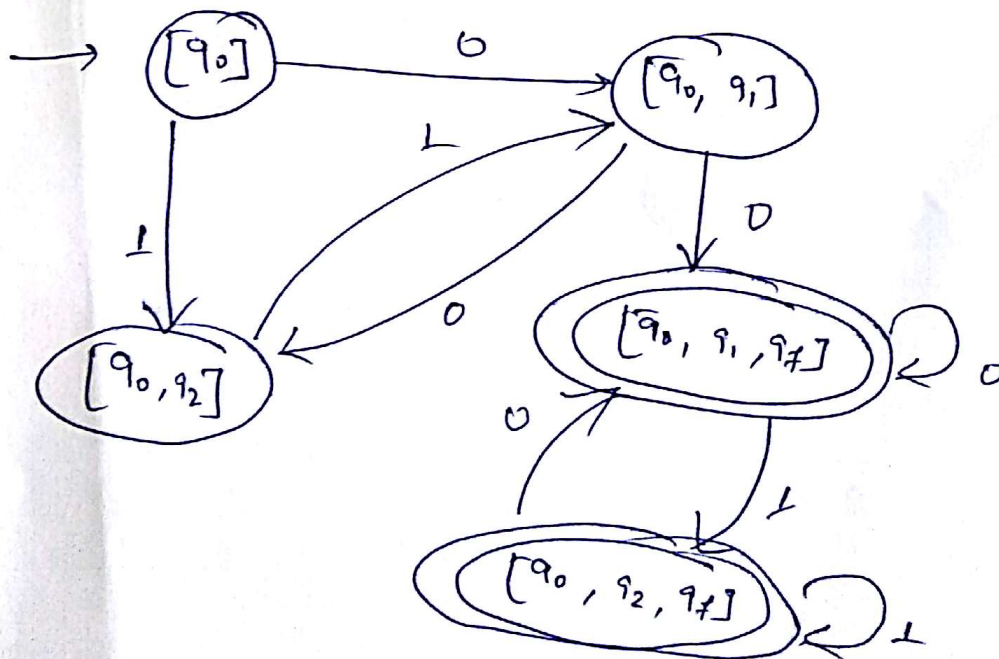
TOC

④

$[q_0], [q_0, q_1], [q_0, q_2], [q_0, q_1, q_f], [q_0, q_2, q_f]$.

Steps:- final state for the resultant DFA
 $= ([q_0, q_1, q_f], [q_0, q_2, q_f])$

States	Input	
	0	1
$\rightarrow [q_0]$	$[q_0, q_1]$	$[q_0, q_2]$
$[q_0, q_1]$	$[q_0, q_1, q_f]$	$[q_0, q_2]$
$[q_0, q_2]$	$[q_0, q_1]$	$[q_0, q_2, q_f]$
$[q_0, q_1, q_f]$	$[q_0, q_1, q_f]$	$[q_0, q_2, q_f]$
$[q_0, q_2, q_f]$	$[q_0, q_1, q_f]$	$[q_0, q_2, q_f]$



This is the required solution.

Elimination of Λ -moves from Transition system TOC

(5)

Q4) Given the steps required to convert a transition system with Λ -moves into an equivalent transition system without Λ -moves.



Step 1: - The pair of vertex which contains Λ -moves
 (q_0, q_1) & (q_1, q_2) .

Step 2: - First of all we have to replace the Λ -transition from vertex (q_0) to the vertex (q_1)

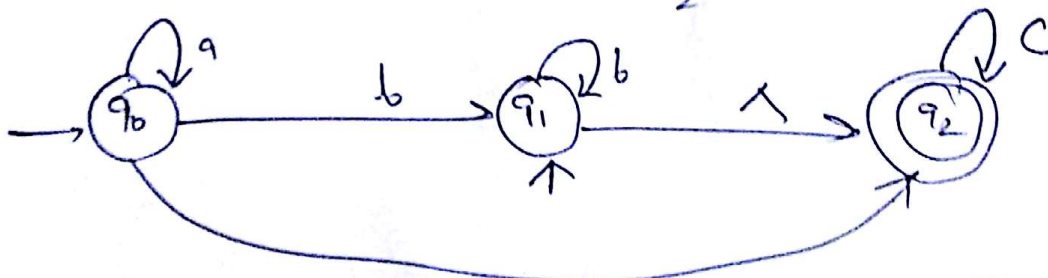
vertex no - 1 = q_0

vertex 2 = q_1

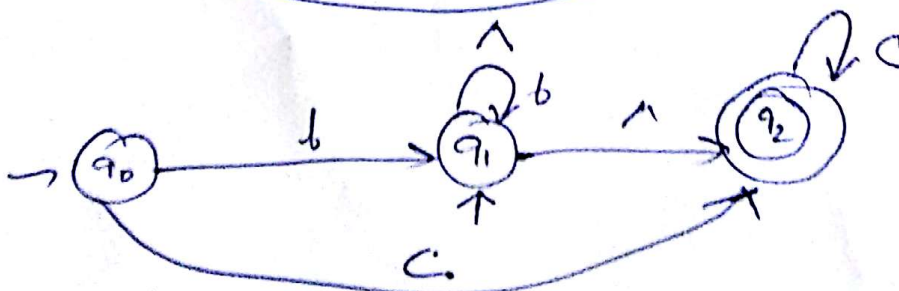
Step 3: - $q_1 \xrightarrow{b} q_1$
 $q_1 \xrightarrow{\Lambda} q_2$

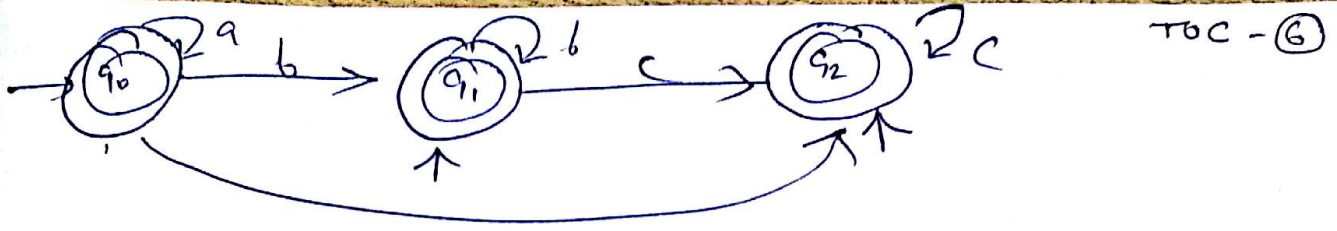
Step 4: - Also, we have to make transition from vertex q_0 .

$q_0 \xrightarrow{b} q_1$
 $q_0 \xrightarrow{\Lambda} q_2$



Steps: -





Solution of question no 4.

Fig: - Desired Transition system with λ (null) moves.